

Math 210A Lecture 3 Notes

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1 The Yoneda Lemma

1.1 Two versions of the Yoneda lemma

Lemma 1.1 (Yoneda). *Let \mathcal{C} be a small category, and let $h^{\mathcal{C}} : \mathcal{C} \rightarrow \text{Fun}(\mathcal{C}^{op}, \text{Set})$ be $h^{\mathcal{C}}(A) = h^A = \text{Hom}_{\mathcal{C}}(\cdot, A)$ and if $f : A \rightarrow B$, then $h^{\mathcal{C}}(f)_X : \text{Hom}(X, A) \rightarrow \text{Hom}(X, B)$ sends $(g : X \rightarrow A) \mapsto (f \circ g : X \rightarrow B)$. Then $h^{\mathcal{C}}$ is fully faithful.*

Proof. To show that $h^{\mathcal{C}}$ is faithful, let $f, g : A \rightarrow B$, and suppose that $h^{\mathcal{C}}(f) = h^{\mathcal{C}}(g)$. Then $h^{\mathcal{C}}(f)_A, h^{\mathcal{C}}(g)_A : \text{Hom}(A, A) \rightarrow \text{Hom}(A, B)$ maps $1_A \mapsto f \circ 1_A = f$ and $1_A \mapsto g \circ 1_A = g$. So $f = g$.

To show that $h^{\mathcal{C}}$ is full, let $\{\eta_X\} : h^A \rightarrow h^B$. We claim that $h^{\mathcal{C}}(\eta_A(1_A)) = \eta$.

$$\begin{array}{ccc} h^A(A) & \xrightarrow{\eta_A} & h^B(A) \\ \downarrow h^A(f) & & \downarrow h^B(f) \\ h^A(C) & \xrightarrow{\eta_C} & h^B(C) \end{array}$$

This is

$$\begin{array}{ccc} \text{Hom}(A, A) & \xrightarrow{\eta_A} & \text{Hom}(B, B) \\ \downarrow h^A(f) & & \downarrow h^B(f) \\ \text{Hom}(C, A) & \xrightarrow{\eta_C} & \text{Hom}(C, B). \end{array}$$

Since this diagram commutes, $\eta_C \circ h^A(f) = h^B(f) \circ \eta_A$. So they are equal on evaluation on an element. Then $\eta_C \circ h^A(f)[1_A] = h^B(f) \circ \eta_A[1_A]$, so $\eta_C[f] = \eta_A[1_A] \circ f$. In particular, $\eta = h^{\mathcal{C}}(\eta_A[1_A])$. \square

Lemma 1.2 (Yoneda, strengthened). *Let \mathcal{C} be a small category, let $h^{\mathcal{C}} : \mathcal{C} \rightarrow \text{Fun}(\mathcal{C}^{op}, \text{Set})$ be the Yoneda embedding, and let $F : \mathcal{C}^{op} \rightarrow \text{Set}$. Then $\text{Nat}(h^A, F)$ is in bijection with $F(A)$.*

Proof. Define $\Phi : \text{Nat}(h^A, F) \rightarrow F(A)$ given by $\eta_A : h^A(A) \rightarrow F(A)$, which sends $1_A \mapsto \eta_A(1_A)$. Define $\Psi : F(A) \rightarrow \text{Nat}(h^A, F)$. Then, for $x \in F(A)$, $\Psi(x)_B : h^A(B) = \text{Hom}(B, A) \rightarrow F(B)$ is $\Psi(x) = \text{ev}_x \circ F$.

We claim that $\Phi \circ \Psi$ is the identity on $F(A)$. Let $x \in F(A)$. Then $\Phi(\Psi(x)) = \Phi(\text{ev}_x \circ F) = \text{ev}_x \circ 1_{F(A)} = x$. $(\Psi \circ \Phi)(\eta) = \Psi(\eta_A(1_A)) = \text{ev}_{\eta_A(1_A)} \circ F$. Let $f : B \rightarrow A$. Then

$$\begin{array}{ccc} \text{Hom}(A, A) & \xrightarrow{\eta_A} & F(A) \\ \downarrow h^A(f) & & \downarrow F(f) \\ \text{Hom}(B, A) & \xrightarrow{\eta_B} & F(B). \end{array}$$

So $F(f) \circ \eta_A = \eta_B \circ h^A(f)$, which means $F(f) \circ \eta_A(1_A) = \eta_B \circ h^A(f)(1_A)$. The left hand side is $\Phi \circ \Phi(\eta)_B[f]$, and the right hand side is $\eta_B(f)$. Therefore, $\Psi \circ \Phi(\eta) = \eta$. \square

This form of the Yoneda lemma implies the previous version.

Corollary 1.1 (Yoneda lemma). *Let \mathcal{C} be a small category, and let $h^{\mathcal{C}} : \mathcal{C} \rightarrow \text{Fun}(\mathcal{C}^{op}, \text{Set})$. Then $h^{\mathcal{C}}$ is fully faithful.*

Proof. Let $B \in \text{Obj}(\mathcal{C})$. Consider $F = h^B - \text{Hom}_{\mathcal{C}}(\cdot, B)$. Then $\text{Nat}(h^A, h^B)$ is in bijection (via F) with $h^B(A) = \text{Hom}_{\mathcal{C}}(A, B)$. \square

1.2 Partially ordered sets

Definition 1.1. A **partially ordered set (poset)** is a set S with a relation \leq on S such that

1. $x \leq x$ for all $x \in S$,
2. if $x \leq y$ and $y \leq x$, then $x = y$,
3. if $x \leq y$ and $y \leq z$, then $x \leq z$.

We can turn a poset into a category. Let $\text{Obj}(\mathcal{C}_S) = S$ and

$$\text{Hom}_{\mathcal{C}_S}(X, Y) = \begin{cases} \{\text{unique morphism}\} & x \leq y \\ \emptyset & \text{otherwise.} \end{cases}$$